### The Power of Abstraction

Barbara Liskov November 2009

# Outline

- Inventing abstract data types
- CLU
- Type hierarchy
- What next



### Data Abstraction Prehistory

The Venus machine

#### The Interdata 3





#### Data Abstraction Prehistory

- The Venus machine
- The Venus operating system



#### Data Abstraction Prehistory

- The Venus machine
- The Venus operating system
- Programming methodology



### Programming Methodology

- How should programs be designed?
- How should programs be structured?



 E. W. Dijkstra. Go To Statement Considered Harmful. Cacm, Mar. 1968



 N. Wirth. Program Development by Stepwise Refinement. Cacm, April 1971



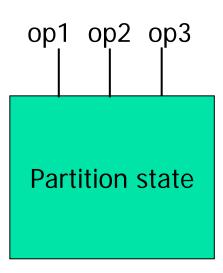
- D. L. Parnas. Information Distribution Aspects of Design Methodology. IFIP Congress, 1971
- "The connections between modules are the assumptions which the modules make about each other."



#### **Partitions**

 B. Liskov. A Design Methodology for Reliable Software Systems. FJCC, Dec. 1972

### Partitions





#### From Partitions to ADTs

How can these ideas be applied to building programs?

# Idea

Connect partitions to data types



#### Meeting in Savanah

- ACM Sigplan-Sigops interface meeting. April 1973. (Sigplan Notices, Sept. 1973)
- Started to work with Steve Zilles



- Extensible Languages
  - S. Schuman and P. Jourrand. Definition Mechanisms in Extensible Programming Languages. AFIPS. 1967
  - R. Balzer. Dataless Programming. FJCC 1967



 O-J. Dahl and C.A.R. Hoare. Hierarchical Program Structures. Structured Programming, Academic Press, 1972



J. H. Morris. Protection in Programming Languages. Cacm. Jan. 1973



 W. Wulf and M. Shaw. Global Variable Considered Harmful. Sigplan Notices. Feb. 1973.



#### **Abstract Data Types**

 B. Liskov and S. Zilles. Programming with Abstract Data Types. ACM Sigplan Conference on Very High Level Languages. April 1974



#### What that paper proposed

- Abstract data types
  - A set of operations
  - And a set of objects
  - The operations provide the only way to use the objects



#### What that paper proposed

- Abstract data types
  - Clusters with encapsulation
- Polymorphism
- Static type checking (we hoped)
- Exception handling



#### From ADTs to CLU

- Participants
  - Russ Atkinson
  - Craig Schaffert
  - Alan Snyder





- Communicating to programmers
- Do ADTs work in practice?
- Getting a precise definition
- Achieving reasonable performance



#### Language Design

- Goals
  - Expressive power, simplicity, performance, ease of use
  - Minimality
  - Uniformity
  - Safety



#### Language Design

- Restrictions
  - No concurrency
  - No go tos
  - No inheritance



#### Some Assumptions/Decisions

- Heap-based with garbage collection!
- No block structure!
- Separate compilation
- Static type checking



#### **CLU Mechanisms**

- Clusters
- Polymorphism
- Exception handling
- Iterators

IntSet = cluster is create, insert, delete, isIn, ...

end IntSet

IntSet = cluster is create, insert, delete, ... end IntSet

```
IntSet s := IntSet$create()
IntSet$insert(s, 3)
```

IntSet = cluster is create, insert, delete, ...

rep = array[int]

```
IntSet = cluster is create, insert, delete, ...
rep = array[int]

create = proc () returns (cvt)
    return (rep$create())
    end create
```

#### Polymorphism

```
Set = cluster[T: type] is create, insert, ... end Set
```

```
Set[int] s := Set[int]$create()
Set[int]$insert(s, 3)
```

# -

#### Polymorphism

```
Set = cluster[T: type] is create, insert, ...
  where T has equal: proctype(T, T)
  returns (bool)
```

#### Polymorphism

```
Set = cluster[T: type] is create, insert, ...
where T has equal: proctype(T, T)
  returns (bool)

rep = array[T]

insert = proc (x: cvt, e: T)
  ... if e = x[i] then ...
```



- J. Goodenough. Exception Handling: Issues and a Proposed Notation. Cacm, Dec. 1975
  - Termination vs. resumption
  - How to specify handlers

```
choose = proc (x: cvt) returns (T)
signals (empty)
if rep$size() = 0 then signal empty
...
```

```
choose = proc (x: cvt) returns (T)
signals (empty)
if rep$size() = 0 then signal empty
...
set[T]$ choose(s)
except when empty: ...
```



- Handling
- Propagating
- Shouldn't happen
  - The failure exception
- Principles
  - Accurate interfaces
  - Avoid useless code



## Iterators

For all x in C do S



### **Iterators**

- For all x in C do S
  - Destroy the collection?
  - Complicate the abstraction?



### Visit to CMU

- Bill Wulf and Mary Shaw, Alphard
- Generators

### **Iterators**

```
sum: int := 0
for e: int in Set[int]$members(s) do
  sum := sum + e
  end
```

# It

### **Iterators**

```
Set = cluster[T] is create, ..., members, ...
rep = array[T]

members = iter (x: cvt) yields (T)
for z: T in rep$elements(x) do
    yield (z) end
```



### After CLU

- Argus and distributed computing
- Type Hierarchy



# The Landscape

- Inheritance was used for:
  - Implementation
  - Type hierarchy



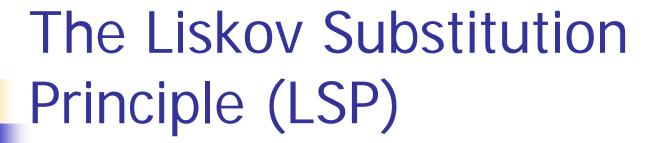
# Implementation Inheritance

Violated encapsulation!



# Type hierarchy

- Wasn't well understood
- E.g., stacks vs. queues



 Objects of subtypes should behave like those of supertypes if used via supertype methods

 B. Liskov. Data abstraction and hierarchy. Sigplan notices, May 1988

# 4

## Polymorphism

- where T has ... vs.
- where T subtype of S

Proofs happen at different times!



 Modularity based on abstraction is the way things are done



# Challenges

- New abstraction mechanisms?
- Massively Parallel Computers
- Internet Computer
  - Storage and computation
  - Semantics, reliability, availability, security

# The Power of Abstraction

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